

Graphic Solution of $2 \times n$ and $n \times 2$ Games.

Graphical method is applicable to only those games in which one of the player has only two strategies. This method reduces the $2 \times n$ or $n \times 2$ game to 2×2 size by identifying and then eliminating the dominant strategies and then solve it by the analytical (algebraic or arithmetic) methods. The resulting solution is the solution to the original problem. This method explains the idea of saddle point graphically.

Consider the following $2 \times n$ games:-

$$\begin{array}{c}
 \text{Player B} \rightarrow \\
 \text{Player A} \downarrow \\
 \begin{array}{c}
 x_0 \\
 1-x_0
 \end{array}
 \left[\begin{array}{cccc}
 y_1 & y_2 & \dots & y_n \\
 a_{11} & a_{12} & \dots & a_{1n} \\
 a_{21} & a_{22} & \dots & a_{2n}
 \end{array} \right]
 \end{array}$$

It's assumed that game has no saddle point. Player A has two strategies A_1 and A_2 which he has strategy of playing with prob. x_0 & $1-x_0$, where $0 \leq x_0 \leq 1$ $\#$

Player B has n strategies which he mixes with probabilities y_1, y_2, \dots, y_n where $y_i \geq 0 \forall i=1, \dots, n$ & $\sum_{i=1}^n y_i = 1$

Obj is to determine optimal value of x_0 ~~and~~
 Expected Payoff for A corresponding to the pure strategies of B are given below:

B's pure strategy

A's expected payoff

1

$$a_{11}x + a_{21}(1-x) = (a_{11}-a_{21})x + a_{21}$$

2

$$a_{12}x + a_{22}(1-x) = (a_{12}-a_{22})x + a_{22}$$

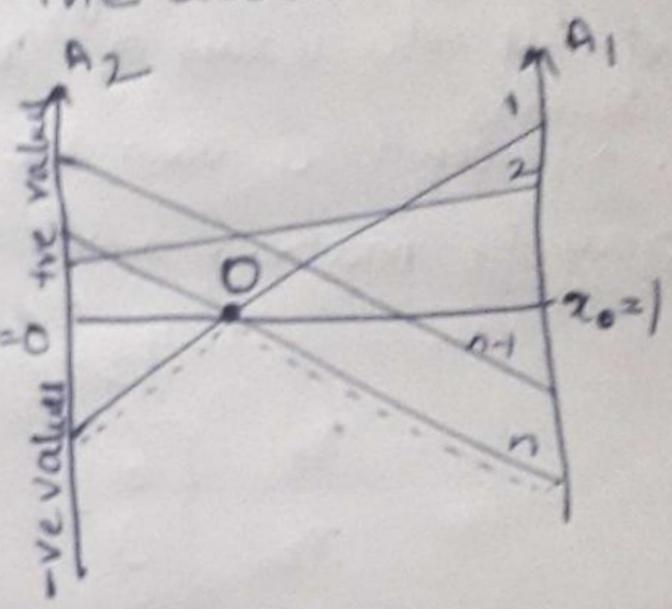
⋮

$$a_{1n}x + a_{2n}(1-x) = (a_{1n}-a_{2n})x + a_{2n}$$

n

∴ A's expected payoffs vary linearly with x . Now according to the maximin criterion for mixed strategies, A should select the value of x which maximizes his minimum expected payoff. This may be done by plotting the above lines as a function of x . As shown graphically.

If A chooses strategy A_1 , $x_1 = 1$ and prob with which he chooses strategy A_2 i.e. $x_2 = 0$



Prob x_1 is represented by the horizontal lines and strategies A_2 and A_1 are represented by two vertical lines unit distance apart.

A's expected payoffs are represented by the slopping lines. A number is allotted to each line corresponding to B's pure strategy. The lower boundary (envelope) of these lines (shown by ---- lines) gives the minimum expected payoff to A as a function of x . A is a maximising player. The highest point on this lower boundary (O) gives the maximum expected payoff to A hence the optimum value of x .

Example: - Solve the given game by graphical method.

$$A \begin{matrix} & y_1 & y_2 & y_3 & y_4 \\ x_1 & 19 & 6 & 7 & 5 \\ x_2 & 7 & 3 & 14 & 6 \\ x_3 & 12 & 8 & 18 & 4 \\ x_4 & 8 & 7 & 13 & -1 \end{matrix}$$

Solⁿ first look for the saddle point using maximin minimax method, which does not exist in the problem.

Now reduce the game by dominance

All the values of column 1 dominate column 2 and all the value of column 3 dominate the " " columns
 \therefore 1 and 3 can be removed from the given payoff matrix and matrix will reduce to (B's minimizing player)

$$\begin{matrix} & y_2 & y_4 \\ x_1 & 6 & 5 \\ x_2 & 3 & 6 \\ x_3 & 8 & 4 \\ x_4 & 7 & -1 \end{matrix}$$

In this matrix row 3 dominates row 4 \therefore row 4 can be removed and reduced matrix is

$$A \begin{matrix} & y_2 & y_4 \\ x_1 & 6 & 5 \\ x_2 & 3 & 6 \\ x_3 & 8 & 4 \end{matrix}$$

Let A_1, A_2, A_3 be the strategies which A mixes with probability x_1, x_2 & x_3 and B mixes

strategy B_2 & B_4 with probability y_2 & y_4

If B adopts B_2 $y_2=1$ & $y_4=0$.

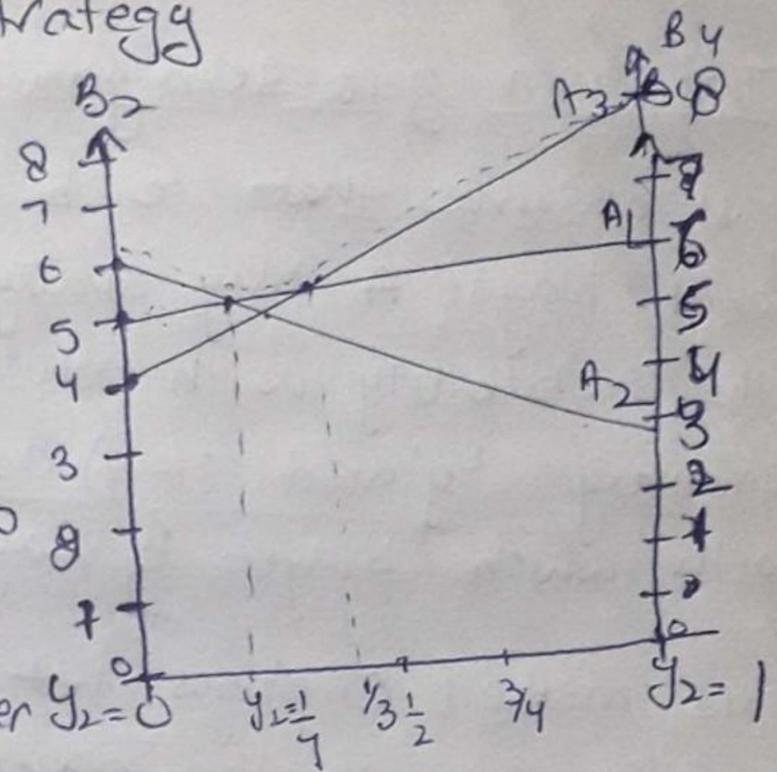
Now we use graphical method to solve the game. B's expected payoff corrⁿ to A's pure strategy are given below

$$\begin{matrix} A_1 & 6y_2 + 5(1-y_2) = y_2 + 5 \\ A_2 & 3y_2 + 6(1-y_2) = -3y_2 + 6 \\ A_3 & 8y_2 + 4(1-y_2) = 4y_2 + 4 \end{matrix}$$

The three lines are straight line & are func of y_2
 draw B_2 & B_4 parallel to each other and one unit apart and mark a scale on each of them.

To represent's A's first strategy corresponds to $y_2 + 5$

when $y_2 = 0$ mark 5 on B_2
 when $y_2 = 1$ mark 6 on B_4
 and join the two points and similarly plot other two strategies also



since B is a minimizing player $y_2 = 0$ bound the figure from above

as shown select the lowest point from the selected bound which is the intersection point of strategy A_1 & A_2 \therefore the above 3×2 game reduces to 2×2 game which can be easily solved by arithmetic method here $y_2 = \frac{1}{4}$

	y_2	y_4	
x_1	6	5	3
x_2	3	6	1
	$\frac{1}{4}$	$\frac{3}{4}$	

\therefore optimal strategy are
 $A = (\frac{3}{4}, \frac{1}{4}, 0, 0)$ $B = (0, \frac{1}{4}, 0, \frac{3}{4})$

$$V = 6 \times \frac{1}{4} + 6 \times \frac{3}{4} \times \frac{1}{4} + 5 \times \frac{3}{4} \times \frac{3}{4} + 3 \times \frac{1}{4} \times \frac{1}{4} + 6 \times \frac{1}{4} \times \frac{3}{4}$$

$$= \frac{18}{16} + \frac{45}{16} + \frac{9}{16} + \frac{18}{16} = \frac{84}{16} = \frac{21}{4} = 5 \frac{1}{4}$$

Solve the following 2×5 game by graphical method:

TABLE 9.132

		Player B				
		1	2	3	4	5
Player A	x_1	-5	5	0	-1	8
	$x_2 = 1 - x_1$	8	-4	-1	6	-5

[P.U.B. Com. 2006; B.B.A.]

The first step is to look for a saddle point. It does not exist in the present problem. The second step is to see if the game can be reduced by dominance. In the present problem, the matrix cannot be reduced by dominance. So, let us solve the matrix by graphical method. A's expected payoffs corresponding to B's pure strategies are

<i>B's pure strategies</i>	<i>A's expected payoffs</i>
1	$-5x_1 + 8(1 - x_1) = -13x_1 + 8$
2	$5x_1 - 4(1 - x_1) = 9x_1 - 4$
3	$0x_1 - 1(1 - x_1) = x_1 - 1$
4	$-1x_1 + 6(1 - x_1) = -7x_1 + 6$
5	$8x_1 - 5(1 - x_1) = 13x_1 - 5$

These five lines can be plotted as functions of x_1 as follows:

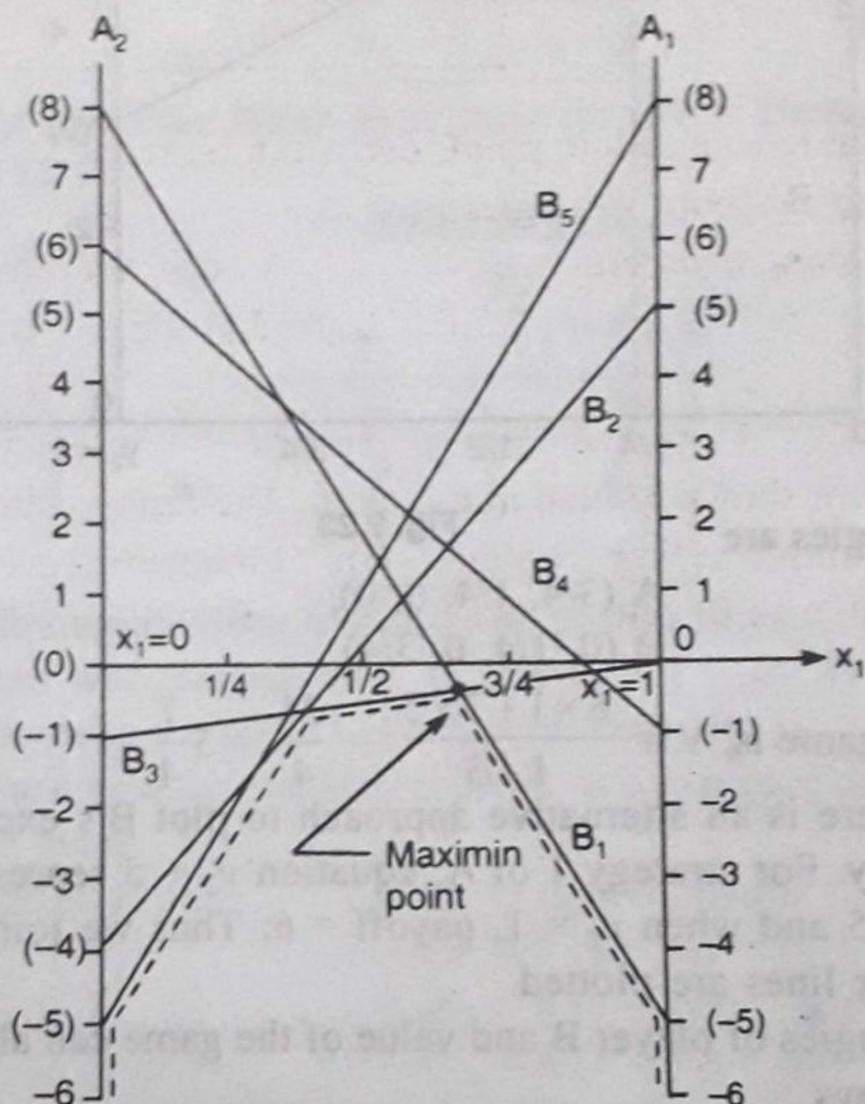


Fig. 9.24

Draw two parallel lines A_1 and A_2 one unit apart and mark a scale on each of them (Fig. 9.24). To represent B's first strategy, join mark -5 on A_1 with mark 8 on A_2 ; to represent B's second strategy, join mark 5 on A_1 with mark -4 on A_2 and so on (for the remaining three strategies) and bound the figure from below.

Since player A wishes to maximize his minimum expected payoff, the two lines which intersect at the highest point of the lower bound show the two courses of action B should choose in his best strategy, which are B_1 and B_3 . We can thus immediately reduce the 2×5 game to 2×2 game which can be easily solved, say, by arithmetic method. The resulting 2×2 game is shown in table 9.133.

TABLE 9.133

		<i>B</i>			
		1	3		
<i>A</i>	1	-5	0	9	9/14
	2	8	-1	5	5/14
		1	13		
		1/14	13/14		

∴ Optimum strategies are

$$A \left(\frac{9}{14}, \frac{5}{14} \right), B \left(\frac{1}{14}, 0, \frac{13}{14}, 0, 0 \right),$$

$$\text{value of the game is, } V = \frac{-5 \times 1 + 0 \times 13}{1 + 13} = \frac{-5}{14}$$

EXAMPLE 9.19-3.3

Solve the following game:

TABLE 9.134

		B		
		y_1	y_2	y_3
A	x_1	6	4	3
	$x_2 = 1 - x_1$	2	4	8

Solution

This game does not have a saddle point, nor it can be reduced by dominance. A's expected payoffs corresponding to B's pure strategies are

B's pure strategies

A's expected payoffs

1

$$6x_1 + 2(1 - x_1) = 4x_1 + 2$$

2

$$4x_1 + 4(1 - x_1) = 4$$

3

$$3x_1 + 8(1 - x_1) = -5x_1 + 8$$

These three lines are shown plotted in Fig. 9.25. The figure is bound from below. There are two maximum points C and D giving the same value of the game.

(i) Point C : Strategies B_1 and B_2 are chosen by B and the following 2×2 game is to be solved :

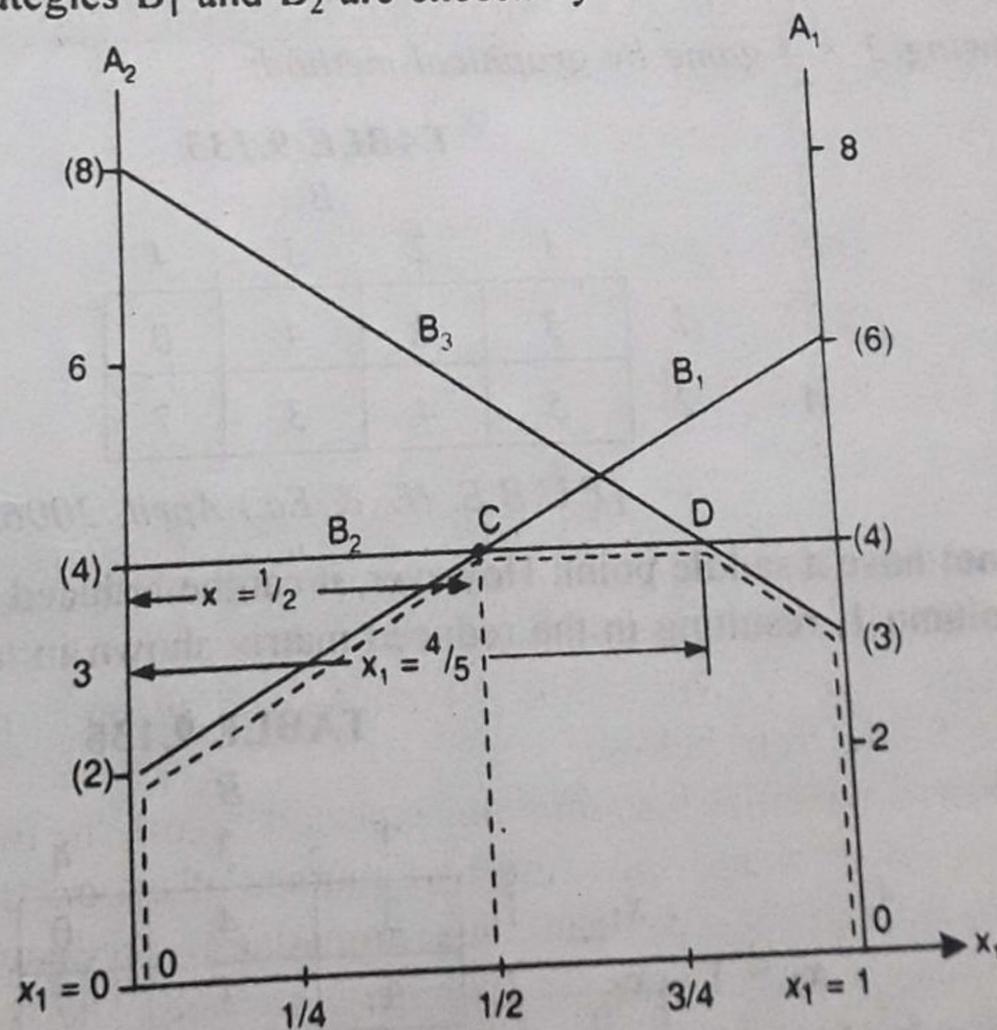


Fig. 9.25

		B			
		1	2		
A	1	6	4	2	$\frac{1}{2}$
	2	2	4	2	$\frac{1}{2}$
		0	4		
		0	1		

∴ Optimal solution is : A $(\frac{1}{2}, \frac{1}{2})$, B (0, 1, 0) ; V = 4.

(ii) Point D : Strategies B₂ and B₃ are mixed by B and the following 2 × 2 game is to be solved:

		B			
		2	3		
A	1	4	3	4	$\frac{4}{5}$
	2	4	8	1	$\frac{1}{5}$
		5	0		
		1	0		

∴ Optimal solution is A $(\frac{4}{5}, \frac{1}{5})$, B (0, 1, 0) ; V = 4.

As evident from Fig. 9.25, any value of x_1 between the points C and D shall be optimal for A. Value of x_1 at points C and D is $\frac{1}{2}$ and $\frac{4}{5}$ respectively. Thus, optimal strategy for A is any pair of $(x_1, x_2 = 1 - x_1)$, where $\frac{1}{2} \leq x_1 \leq \frac{4}{5}$ or $0.5 \leq x_1 \leq 0.8$. Optimal strategy for B is (0, 1, 0) and value of game is 4.

Note: If there happen to be more than two lines passing through the maximin (or minimax) point, then any two lines having opposite signs for their slopes will yield an alternative optimal solution.

EXAMPLE 9.19-3.4

Solve the following 2 × 4 game by graphical method:

TABLE 9.135

		B			
		1	2	3	4
A	1	3	3	4	0
	2	5	4	3	7

Solution

[P.U.B.E. (E. & Ec.) April, 2006, B. Com. Sept., 2005]

This game does not have a saddle point. However, it can be reduced by dominance because column 2 dominates column 1, resulting in the reduced matrix shown in table 9.136.

TABLE 9.136

		B			
		2	3	4	
A	x_1	1	3	4	0
	$x_2 = 1 - x_1$	2	4	3	7

This matrix can be solved by graphical method. A's expected payoffs corresponding to B's pure strategies are

B's pure strategies

A's expected payoffs

2

$$3x_1 + 4(1 - x_1) = -x_1 + 4$$

3

$$4x_1 + 3(1 - x_1) = x_1 + 3$$

4

$$0x_1 + 7(1 - x_1) = -7x_1 + 7$$

These three straight lines can be plotted as functions of x_1 . The method of plotting them has already been described; and they appear as shown in figure 9.26.

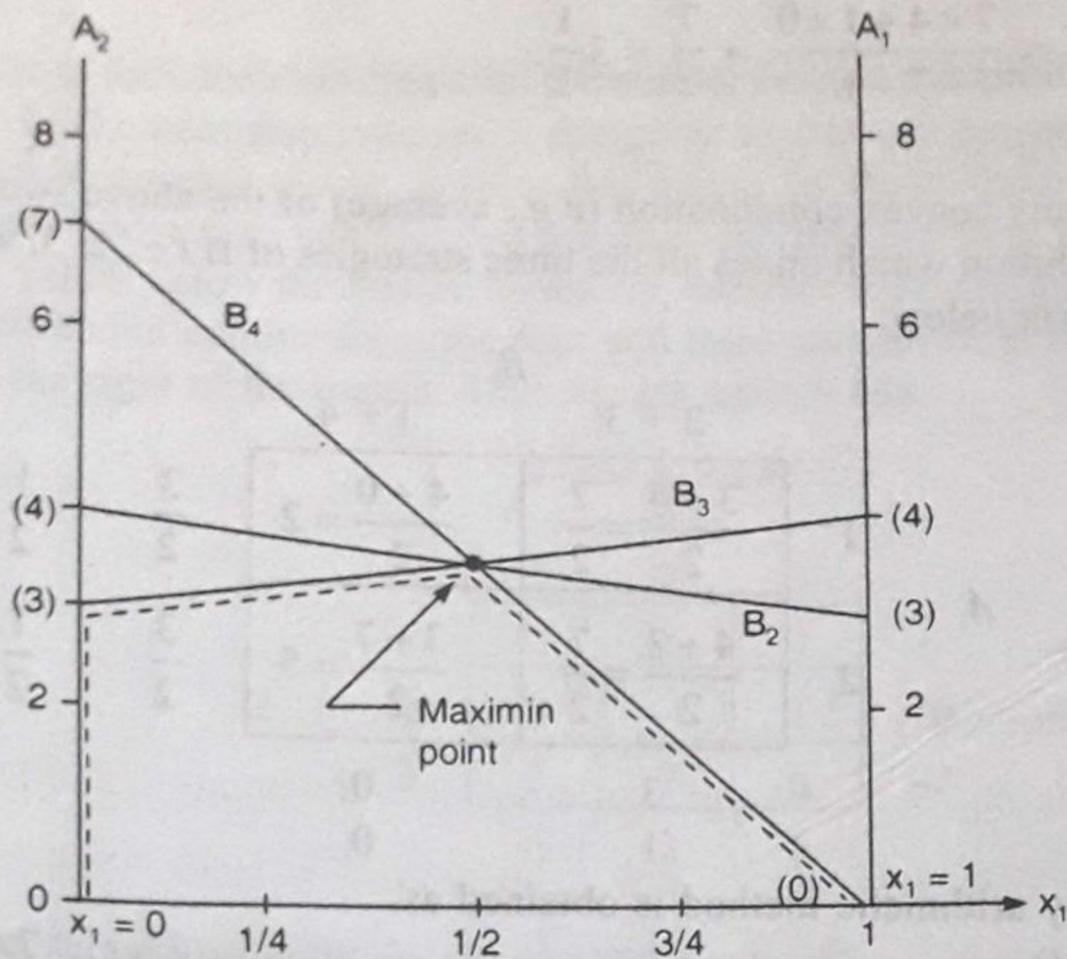


Fig. 9.26

Figure is to be bound from below as shown. All the three lines pass through the maximin point. As mentioned earlier, any two lines having opposite signs for their slopes will yield an alternative optimal solution. This means that the combination of B_2 and B_4 must be excluded as both have the same sign for their slopes. So the game reduces to two 2×2 games which can be easily solved, say, by arithmetic method.

(i)

		<i>B</i>			
		2	3		
<i>A</i>	1	3	4	1	1/2
	2	4	3	1	1/2
		1	1		
		1/2	1/2		

A's optimum strategy : A (1/2, 1/2),

B's optimum strategy : B (0, 1/2, 1/2, 0),

$$\text{game value, } V = \frac{3 \times 1 + 4 \times 1}{1 + 1} = \frac{7}{2} = 3 \frac{1}{2}.$$

(ii)

		<i>B</i>				
		3	4			
<i>A</i>	1	4	0	4	1/2	4/8
	2	3	7	4	1/2	4/8
		7	1			
		7/8	1/8			

A's optimum strategy : A (4/8, 4/8),

B's optimum strategy : B (0, 0, 7/8, 1/8),